

# Advanced Algorithms

Lecture 10: MST (contd.), local search

# Announcements

- HW 2 due tomorrow!
- HW 1 grading, comments (Vivek Gupta)

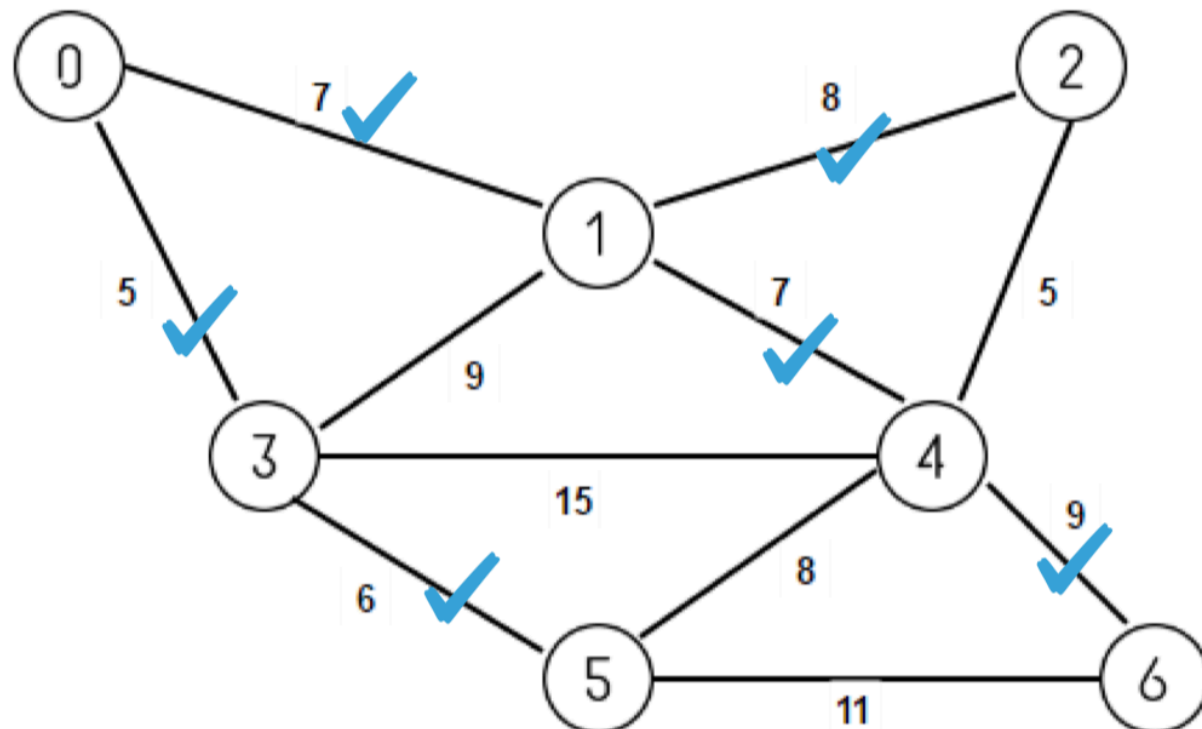
# Greedy algorithms – comments

- Usually “easy” to come up with (we are naturally myopic)
- Usually not optimal – examples, Traveling salesman, set cover, ...
- (Due to this..) analysis is usually tricky

# Example 2: spanning trees

**Problem:** let  $G = (V, E)$  be a (simple, undirected) graph with edge weights  $\{w_e\}$  ( $>0$ ). Pick a subset of the edges, such that (a) all vertices are “connected”, (b) total weight of edges is minimized

(Communication backbone in a network)



# Greedy strategy

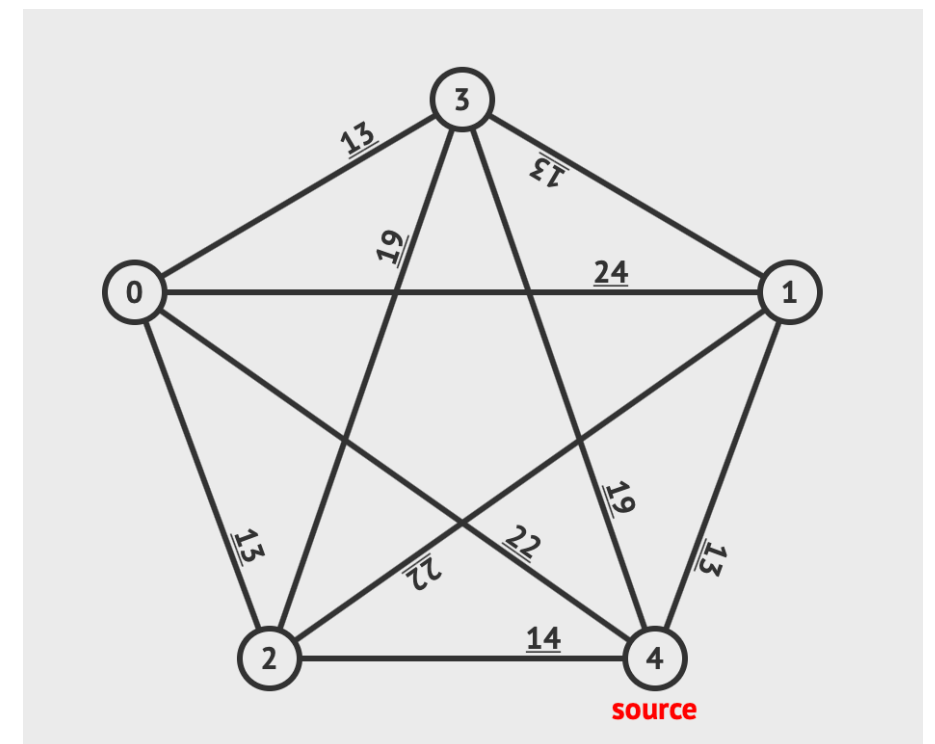
- **Goal:** need to connect all vertices to one another
- Prim: Start with one vertex, add a new vertex to connected set each time
- Kruskal: Add edges one at time, choose min weight edge that isn't "redundant"

**Surprise:** both turn out to be optimal!

# Prim's algorithm

<https://visualgo.net/en/mst>

- start with  $S_1 = \{u\}$
- for  $t = 1, \dots, n-1$ :
  - add least wt edge out of  $S_t$



# Correctness

- **Observation:** at each iteration, we have a *set of connected vertices* —  $S_t$
- Will show: There exists a min spanning tree for the *full graph* that contains all edges chosen so far — **structural assumption**

Inductive proof: assuming there's an MST for the full graph containing edges added until  $t$ , prove that there's an MST for the full graph containing edge added at  $t+1$

# Proof of "opt prefix" property



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Running time

# Minimum spanning tree

- Simple algorithms — analysis slightly tricky
- Common inductive approach for greedy algorithms: show that there's an optimal solution that agrees with all choices so far
- Can be solved in  $O((m + n) \log n)$  time
- Procedure closely related to shortest paths — Dijkstra's algorithm

Local search

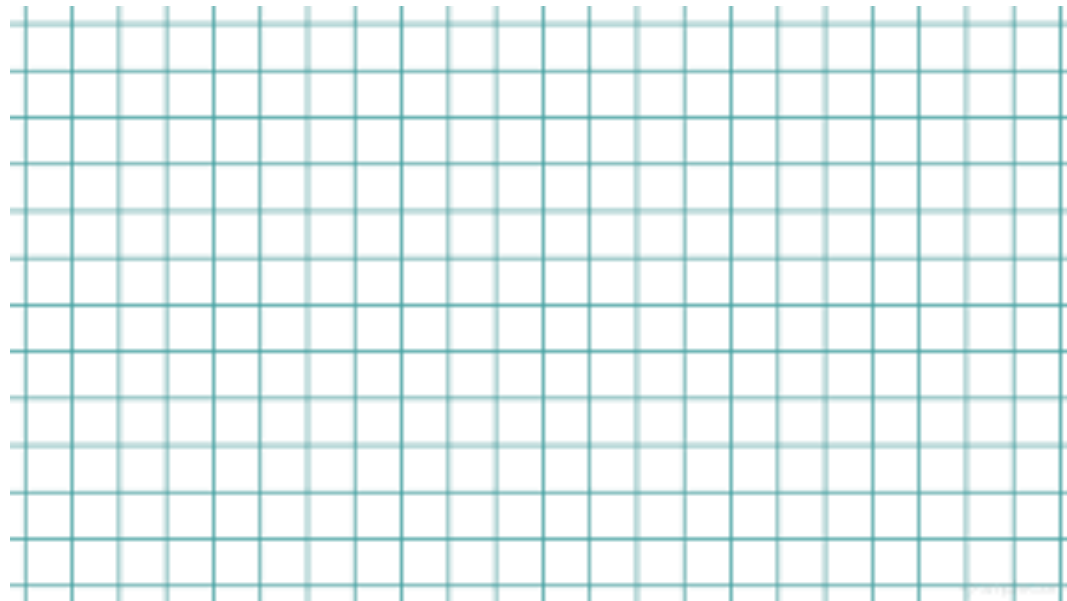
# Main idea

- Start with *any* solution, try improving by moving to “nearby” solution
- Stop if no nearby solution is better

# Classic example – function opt

**Problem:** Let  $f(x)$  be a function defined on domain  $D$ . Find  $\operatorname{argmin}_x f(x)$

# Multi-variate functions



# When is it optimal?

- Any *local optimum* is actually “global” optimum (opt over domain)
- Does this property hold for some natural class?



# Minimizing a convex function

(over all of  $\mathbb{R}^n$ )

# Gradient descent

**(all of modern ML)**

- What is a direction in which function value drops?
- General algorithm:

# Matching problem

**Problem:** suppose we have  $n$  children and  $n$  gifts. Each child has some “happiness value” ( $V_{ij}$ ) for each gift. Find an allocation (one gift per child) so that total happiness is maximized.



Matching – greedy?

# Local search

# Local search

**Claim:** take any solution  $S$  in which swaps do not increase value. Then total happiness of  $S \geq (1/2)$  total happiness of OPT solution

# 2 approximation – proof

**Claim:** take any solution  $S$  in which swaps do not increase value. Then total happiness of  $S \geq (1/2)$  total happiness of OPT solution